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APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
10/661,852	09/12/2003	Mandayam Thondanur Raghunath	YOR920030222US1	8538
24299	7590	02/13/2007		
GEORGE SAI-HALASZ 303 TABER AVENUE PROVIDENCE, RI 02906			EXAMINER TRAN, ELLEN C	
			ART UNIT 2134	PAPER NUMBER

SHORTENED STATUTORY PERIOD OF RESPONSE	MAIL DATE	DELIVERY MODE
3 MONTHS	02/13/2007	PAPER

Please find below and/or attached an Office communication concerning this application or proceeding.

If NO period for reply is specified above, the maximum statutory period will apply and will expire 6 MONTHS from the mailing date of this communication.

Office Action Summary	Application No. 10/661,852	Applicant(s) RAGHUNATH ET AL.	
	Examiner Ellen C. Tran	Art Unit 2134	

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

Period for Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

Status

- 1) ☒ Responsive to communication(s) filed on 12 September 2003.
- 2a) ☐ This action is **FINAL**. 2b) ☒ This action is non-final.
- 3) ☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

Disposition of Claims

- 4) ☒ Claim(s) 1-20 is/are pending in the application.
- 4a) Of the above claim(s) _____ is/are withdrawn from consideration.
- 5) ☐ Claim(s) _____ is/are allowed.
- 6) ☒ Claim(s) 1-20 is/are rejected.
- 7) ☐ Claim(s) _____ is/are objected to.
- 8) ☐ Claim(s) _____ are subject to restriction and/or election requirement.

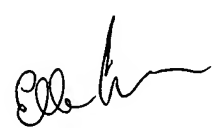
Application Papers

- 9) ☐ The specification is objected to by the Examiner.
- 10) ☐ The drawing(s) filed on _____ is/are: a) ☐ accepted or b) ☐ objected to by the Examiner.
Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).
Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).
- 11) ☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

Priority under 35 U.S.C. § 119

- 12) ☐ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).
- a) ☐ All b) ☐ Some * c) ☐ None of:
1. ☐ Certified copies of the priority documents have been received.
2. ☐ Certified copies of the priority documents have been received in Application No. _____.
3. ☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

* See the attached detailed Office action for a list of the certified copies not received.



Attachment(s)

- | | |
|--|---|
| 1) <input checked="" type="checkbox"/> Notice of References Cited (PTO-892) | 4) <input type="checkbox"/> Interview Summary (PTO-413)
Paper No(s)/Mail Date. _____ |
| 2) <input type="checkbox"/> Notice of Draftsperson's Patent Drawing Review (PTO-948) | 5) <input type="checkbox"/> Notice of Informal Patent Application |
| 3) <input checked="" type="checkbox"/> Information Disclosure Statement(s) (PTO/SB/08)
Paper No(s)/Mail Date <u>8 June 2004</u> . | 6) <input type="checkbox"/> Other: _____ |

DETAILED ACTION

1. This action is responsive to: an original application filed on 12 September 2003.
2. Claims 1-20 are pending; claims 1, 4, 10, and 20 are independent claims.
3. The IDS submitted 8 June 2004 has been considered.

Claim Rejections - 35 USC § 101

4. 35 U.S.C. 101 reads as follows:

Whoever invents or discovers any new and useful process, machine, manufacture, or composition of matter, or any new and useful improvement thereof, may obtain a patent therefor, subject to the conditions and requirements of this title.

5. Claim 20 is rejected under 35 U.S.C. 101 because the claimed invention is directed to non-statutory subject matter. Claim 20 is directed to carrier wave signal, which is not patentable subject matter.
6. To expedite a complete examination of the instant application the claims rejected under 35 U.S.C. 101 (nonstatutory) are further rejected as set forth below in anticipation of applicant amending these claims to place them within the four statutory categories of invention.

Claim Rejections - 35 USC § 103

7. The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negated by the manner in which the invention was made.

8. **Claims 1, 4-8, 10-15, and 18-20**, are rejected under 35 U.S.C. 103(a) as being unpatentable over Kocher et al. U.S. Patent No. 6,304,658 (hereinafter '658).

As to independent claim 1, “a first copy of a shared secret key” is shown in ‘658 col. 5, lines 43-46 “Assume a user wishes to authenticate herself to a server using an n-bit secret key, K, known to both the server and the user's cryptographic token, but not known to attackers”;

“a first standard certificate, wherein the first standard certificate is being used in responding to a challenge of the door” is disclosed in ‘658 col. 11, lines 9-21 “Using techniques of the background art, Alice and Bob can use their certificates to establish a secure communication channel. They first exchange certificates (C_{Alice} and C_{Bob}). Each verifies that the other's certificate is acceptable (e.g., properly formatted, properly signed by a trusted CA, not expired, not revoked, etc.). Because this protocol will assume that p and g are constants, they also check that the certificate's p and g match the expected values”;

“and means for communicating with the door, wherein the door possesses a second copy of the shared secret key, and wherein the door adapted to validate identicalness of the first and the second copies of the shared secret key” is taught ‘658 col. 21, lines 22-34 and col. 5, lines 51-59, note the ‘wireless keys’ are a means of communicating with the door and the second copy is equivalent to the servers copy of K;

“and wherein the door further adapted to issue the challenge to the computing device” is taught in col. 5, lines 51-55 “Typically the server sends a unique, unpredictable challenge value R to the user's token, which computes the value $A = H(R.\text{parallel}.K)$, where “.parallel.” denotes concatenation and H is a one-way cryptographic hash function such as SHA”; the following is not explicitly taught in ‘658: **As to independent claim 1, “A portable computing device for opening a door, comprising: a memory, wherein a content of the memory comprises:”** however ‘658 teaches “Leak-resistant algorithms, protocols, and devices

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may be used in virtually any application requiring cryptographic security and secure key management, including without limitation: smartcards, electronic cash, electronic payments, funds transfer, remote access, timestamping, certification, certificate validation, secure e-mail, secure facsimile, telecommunications security (voice and data), computer networks, radio and satellite communications, infrared communications, access control, door locks, wireless keys, biometric devices, automobile ignition locks, copy protection devices, payment systems, systems for controlling the use and payment of copyrighted information, and point of sale terminals” in col. 21, lines 22-34. Note the portable devices could be interpreted equivalent to smartcards. The invention is described with respect to a server that the smartcard authenticates itself with, the server system could be interpreted to be equivalent to door locks.

It would have been obvious to one of ordinary skill in the art at the time of a method for securing cryptographic devices against attacks involving external monitoring and analysis taught in ‘658 to include a means to utilize these devices for opening a door. One of ordinary skill in the art would have been motivated to perform such a modification because to provide secure key management see ‘500 (col. 2, lines 4 et seq.) “Most cryptosystems require secure key management. In public-key based security systems, private keys must be protected so that attackers cannot use the keys to forge digital signatures, modify data, or decrypt sensitive information. Systems employing symmetric cryptography similarly require that keys be kept secret. Well-designed cryptographic algorithms and protocols should prevent attackers who eavesdrop on communications from breaking systems. However, cryptographic algorithms and protocols traditionally require that tamper-resistant hardware or other implementation-specific measures prevent attackers from accessing or finding the keys”.

As to independent claim 4, “providing a portable computing device, wherein the computing device is equipped with a memory, and the memory holds a first copy of the shared secret key” is shown in ‘658 col. 5, lines 43-46 “Assume a user wishes to authenticate herself to a server using an n-bit secret key, K, known to both the server and the user's cryptographic token, but not known to attackers”;

“and a first standard certificate, wherein the computing device is adapted for performing operations with shared secret keys and standard certificates,” is disclosed in ‘658 col. 11, lines 9-21 “Using techniques of the background art, Alice and Bob can use their certificates to establish a secure communication channel. They first exchange certificates (C_{Alice} and C_{Bob}). Each verifies that the other's certificate is acceptable (e.g., properly formatted, properly signed by a trusted CA, not expired, not revoked, etc.). Because this protocol will assume that p and g are constants, they also check that the certificate's p and g match the expected values”;

“and wherein the computing device is also having means for communicating with the door communicating by the computing device to the door a device identifier; issuing a challenge by the door to the computing device, wherein the challenge is issued only on randomly selected occasions” ” is taught ‘658 col. 21, lines 22-34 and col. 5, lines 51-59, note the ‘wireless keys’ are a means of communicating with the door and the second copy is equivalent to the servers copy of K;

“responding to the challenge by the computing device by demonstrating possession of a private key part of the first standard certificate; responding by the door with a door identifier and with a message, wherein the message is encrypted with a second copy of the

shared secret key, and wherein using the second copy of the shared secret key for encrypting the message resulted from recognizing the device identifier communicated by the computing device” is shown in col. 11, lines 31-45, note after the two parties authenticate each other using certificates they derive a symmetric key, using the device identifier, i.e. public key for future communications;

“responding by the computing device with a signal attesting decryption of the message, wherein the message has been decrypted in the computing device by the first copy of the shared secret key, and wherein using the first copy of the shared secret key for decrypting the message resulted from recognizing the door identifier transmitted by the door; and unlocking the door upon recognizing validity of the signal attesting decryption of the message” is disclosed in ‘658 col. 11, lines 9-21 ,note the signal is the verification; the following is not explicitly taught in ‘658: **“A method for secure unlocking of a door based on a shared secret key, comprising the steps of:”** however ‘658 teaches “Leak-resistant algorithms, protocols, and devices may be used in virtually any application requiring cryptographic security and secure key management, including without limitation: smartcards, electronic cash, electronic payments, funds transfer, remote access, timestamping, certification, certificate validation, secure e-mail, secure facsimile, telecommunications security (voice and data), computer networks, radio and satellite communications, infrared communications, access control, door locks, wireless keys, biometric devices, automobile ignition locks, copy protection devices, payment systems, systems for controlling the use and payment of copyrighted information, and point of sale terminals” in col. 21, lines 22-34. Note the portable devices could be interpreted equivalent to smartcards. The invention is described with respect to a server that

the smartcard authenticates itself with, the server system could be interpreted to be equivalent to door locks.

It would have been obvious to one of ordinary skill in the art at the time of a method for securing cryptographic devices against attacks involving external monitoring and analysis taught in '658 to include a means to utilize these devices for opening a door. One of ordinary skill in the art would have been motivated to perform such a modification because to provide secure key management see '500 (col. 2, lines 4 et seq.) "Most cryptosystems require secure key management. In public-key based security systems, private keys must be protected so that attackers cannot use the keys to forge digital signatures, modify data, or decrypt sensitive information. Systems employing symmetric cryptography similarly require that keys be kept secret. Well-designed cryptographic algorithms and protocols should prevent attackers who eavesdrop on communications from breaking systems. However, cryptographic algorithms and protocols traditionally require that tamper-resistant hardware or other implementation-specific measures prevent attackers from accessing or finding the keys".

As to dependent claim 5, "wherein the device identifier is a hash code of the first standard certificate" is taught in '658 col. 5, lines 51-55.

As to dependent claim 6, "wherein the door identifier is a simple identifier and it is sent without encryption" is shown in '658 col. 5, lines 51-55 the door identifier is the random number which is not encrypted.

As to dependent claim 7, "wherein the door has a second standard certificate, and the door identifier is a hash code of the second standard certificate" is '658 disclosed in col. 11, lines 9-21.

As to dependent claim 8, “wherein the shared secret key is generated by the door and communicated with the computing device in private using a public key part of the first standard certificate” is disclosed in ‘658 col. 11, lines 9-21.

As to independent claim 10, “each computing device equipped with a memory, wherein any one of the computing devices holds in its memory a unique first standard certificate and wherein the any one computing device further holds in its memory door identifiers for all those doors out of the first plurality of doors that the any one computing device is permitted to open” is disclosed in ‘658 col. 11, lines 9-21

“and wherein each of the door identifier is uniquely linked to a first copy of a shared secret key” is shown in ‘658 col. 5, lines 43-46

“wherein any one of the doors possesses a matching information for each one of those computing devices out of the second plurality of computing devices that are permitted to open the any one door, wherein the matching information comprises a device identifier, wherein the device identifier is linked to a public key part of the unique first standard certificate and to a second copy of the shared secret key” is taught ‘658 col. 21, lines 22-34 and col. 5, lines 51-59;

“and wherein the first plurality of doors and the second plurality of computing devices have means for communicating between any device and any door, and wherein the any one door is adapted to recognize the device identifier, and further adapted to use the matching information to validate identicalness of the first and the second copies of the shared secret key, and to issue a challenge to the unique first standard certificate using the public key part of the unique first standard certificate” is taught in col. 5, lines 51-67;

the following is not explicitly taught in '658: **“A security system for controlling access, comprising a first plurality of doors and a second plurality of portable computing devices for opening doors”** however '658 teaches how portable devices, i.e. smartcards can be used with respect to a server that the smartcard authenticates itself with, the server system could be interpreted to be equivalent to door locks in col. 21, lines 22-34.

It would have been obvious to one of ordinary skill in the art at the time of a method for securing cryptographic devices against attacks involving external monitoring and analysis taught in '658 to include a means to utilize these devices for opening a door. One of ordinary skill in the art would have been motivated to perform such a modification because to provide secure key management see '500 (col. 2, lines 4 et seq.) “Most cryptosystems require secure key management. In public-key based security systems, private keys must be protected so that attackers cannot use the keys to forge digital signatures, modify data, or decrypt sensitive information. Systems employing symmetric cryptography similarly require that keys be kept secret. Well-designed cryptographic algorithms and protocols should prevent attackers who eavesdrop on communications from breaking systems. However, cryptographic algorithms and protocols traditionally require that tamper-resistant hardware or other implementation-specific measures prevent attackers from accessing or finding the keys”.

As to dependent claims 11-13, these claims contain substantially similar subject matter as claims 5-7; therefore they are rejected along similar rationale.

As to dependent claim 14, **“wherein the door identifier is a hash code of the unique second standard certificate”** is taught in '658 col. 5, lines 51-55.

As to dependent claim 15, “wherein the challenge is issued on randomly selected occasions” is shown in ‘658 col. 5, lines 51-59.

As to dependent claim 18, “wherein the challenge by the any one door is successfully responded by demonstrating possession of a private key part of the unique first standard certificate” is disclosed in col. 11, lines 9-21.

As to dependent claim 19, “wherein the any one door is further adapted to generate a shared secret key and communicate the shared key in private by using the public key part of the unique first standard certificate” is taught in ‘658 col. 11, lines 9-21.

As to independent claim 20, this claim is directed to a computer data signal that incorporates the limitation of the method of claim 4; therefore it is rejected along similar rationale.

8. **Claims 2, 3, 9, 16 and 17**, are rejected under 35 U.S.C. 103(a) as being unpatentable over Kocher et al. U.S. Patent No. 6,304,658 (hereinafter ‘658) in view of Kawan et al. US Patent No. 7,039,812 (hereinafter ‘812).

As to dependent claim 2, the following is not specifically taught in ‘658: “**wherein the first standard certificate is having a private key part and the private key part is being encrypted with a first biometric key, wherein the first biometric key belongs to a rightful owner of the computing device**” however ‘812 teaches “In an embodiment of the present invention, the user presents one or more user credentials for enrollment by an authority, such as a bank. The user credentials include, for example, one or more biometric templates for the user's fingerprint(s), face, voice and/or iris and/or one or more digital documents, such as a digital certificate and/or a digital signature for the user and/or one or more paper documents, such as a

passport for the user. The user credential(s), which represent user authentication information, are stored for the user, for example, on a host computer, a local terminal, and/or a user token, such as a smart card, and the stored user credential(s) can be signed with the user's private key" in col. 2, lines 17-29.

It would have been obvious to one of ordinary skill in the art at the time of a method for securing cryptographic devices against attacks involving external monitoring and analysis taught in '658 to include a means to utilize biometrics. One of ordinary skill in the art would have been motivated to perform such a modification because to augment the security provided see '812 (col. 60 et seq.). "Most cryptosystems require secure key management. In public-key based security systems, private keys must be protected so that attackers cannot use the keys to forge digital signatures, modify data, or decrypt sensitive information. Systems employing symmetric cryptography similarly require that keys be kept secret. Well-designed cryptographic algorithms and protocols should prevent attackers who eavesdrop on communications from breaking systems. However, cryptographic algorithms and protocols traditionally require that tamper-resistant hardware or other implementation-specific measures prevent attackers from accessing or finding the keys".

As to dependent claim 3, "further comprising a biometric device, wherein the biometric device is capable of generating a second biometric key, wherein the second biometric key belongs to a user of the computing device, and wherein the second biometric key is used to decrypt the private key part of the first standard certificate" however '812 teaches "In yet a further aspect of user authentication for an embodiment of the present invention, the authority 18 takes the form of a user token, such as a smart card 66. The user

credentials 24, such as fingerprints 28, are stored in the user's smart card 66. The shared secret 42 is also stored on the smart card 66. The user 10 can present his credentials 24 to the smart card 66 in the predefined shared secret sequence 46 and open the smart card 66 for its normal usage. In this aspect, the user's credentials 24, such as biometric templates, digital certificates, and the like, verification parameters, and shared secrets 42 can be signed with the user's private key and stored locally for fraud prevention, such as smart card tampering" in col. 10, lines 26-39. The motivation to combine '658 and '812 is the same as stated above in dependent claim 2.

As to dependent claim 9. "wherein the private key part of the first standard certificate is encrypted with a first biometric key, wherein the first biometric key belongs to a rightful owner of the computing device" however '812 teaches biometric certificates encrypted with private keys in col. 2, lines 17-29;

"and wherein the computing device is provided with a biometric device, and wherein the step of responding to the challenge further comprise the steps of: taking a biometric reading of a user of the computing device; generating a second biometric key using the biometric reading; and decrypting the encrypted private key part of the first standard certificate using the second biometric key, whereby if the first and second biometric keys are identical the decrypting using the second biometric key is successful, and the challenge can be successfully responded" however '812 teaches "In yet a further aspect of user authentication for an embodiment of the present invention, the authority 18 takes the form of a user token, such as a smart card 66. The user credentials 24, such as fingerprints 28, are stored in the user's smart card 66. The shared secret 42 is also stored on the smart card 66. The user 10 can present his credentials 24 to the smart card 66 in the predefined shared secret

sequence 46 and open the smart card 66 for its normal usage. In this aspect, the user's credentials 24, such as biometric templates, digital certificates, and the like, verification parameters, and shared secrets 42 can be signed with the user's private key and stored locally for fraud prevention, such as smart card tampering” in col. 10, lines 26-39. The motivation to combine ‘658 and ‘812 is the same as stated above in dependent claim 2.

As to dependent claim 16, “wherein the unique first standard certificate is having a private key part and the private key part is being encrypted with a first biometric key, wherein the first biometric key belongs to a rightful owner of the computing device” however ‘812 teaches biometric certificates encrypted with private keys in col. 2, lines 17-29. The motivation to combine ‘658 and ‘812 is the same as stated above in dependent claim 2.

As to dependent claim 17, “wherein the any one computing device is further comprising a biometric device, wherein the biometric device is capable of generating a second biometric key, wherein the second biometric key belongs to a user of the any one computing device, and wherein the second biometric key is used to decrypt the private key part of the unique first standard certificate” however ‘812 teaches “In yet a further aspect of user authentication for an embodiment of the present invention, the authority 18 takes the form of a user token, such as a smart card 66. The user credentials 24, such as fingerprints 28, are stored in the user's smart card 66. The shared secret 42 is also stored on the smart card 66. The user 10 can present his credentials 24 to the smart card 66 in the predefined shared secret sequence 46 and open the smart card 66 for its normal usage. In this aspect, the user's credentials 24, such as biometric templates, digital certificates, and the like, verification parameters, and shared secrets 42 can be signed with the user's private key and stored locally for fraud

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prevention, such as smart card tampering” in col. 10, lines 26-39. The motivation to combine ‘658 and ‘812 is the same as stated above in dependent claim 2.

Conclusion

9. Any inquiry concerning this communication or earlier communications from the examiner should be directed to Ellen C Tran whose telephone number is (571) 272-3842. The examiner can normally be reached from 6:00 am to 4:00 pm.

If attempts to reach the examiner by telephone are unsuccessful, the examiner’s supervisor, Kambiz Zand can be reached on (571) 272-3811. The fax phone number for the organization where this application or proceeding is assigned is (571) 273-8300.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see <http://pair-direct.uspto.gov>. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free).



Ellen Tran
Patent Examiner
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10 February 2007